

# Introduction to Quantum Computing



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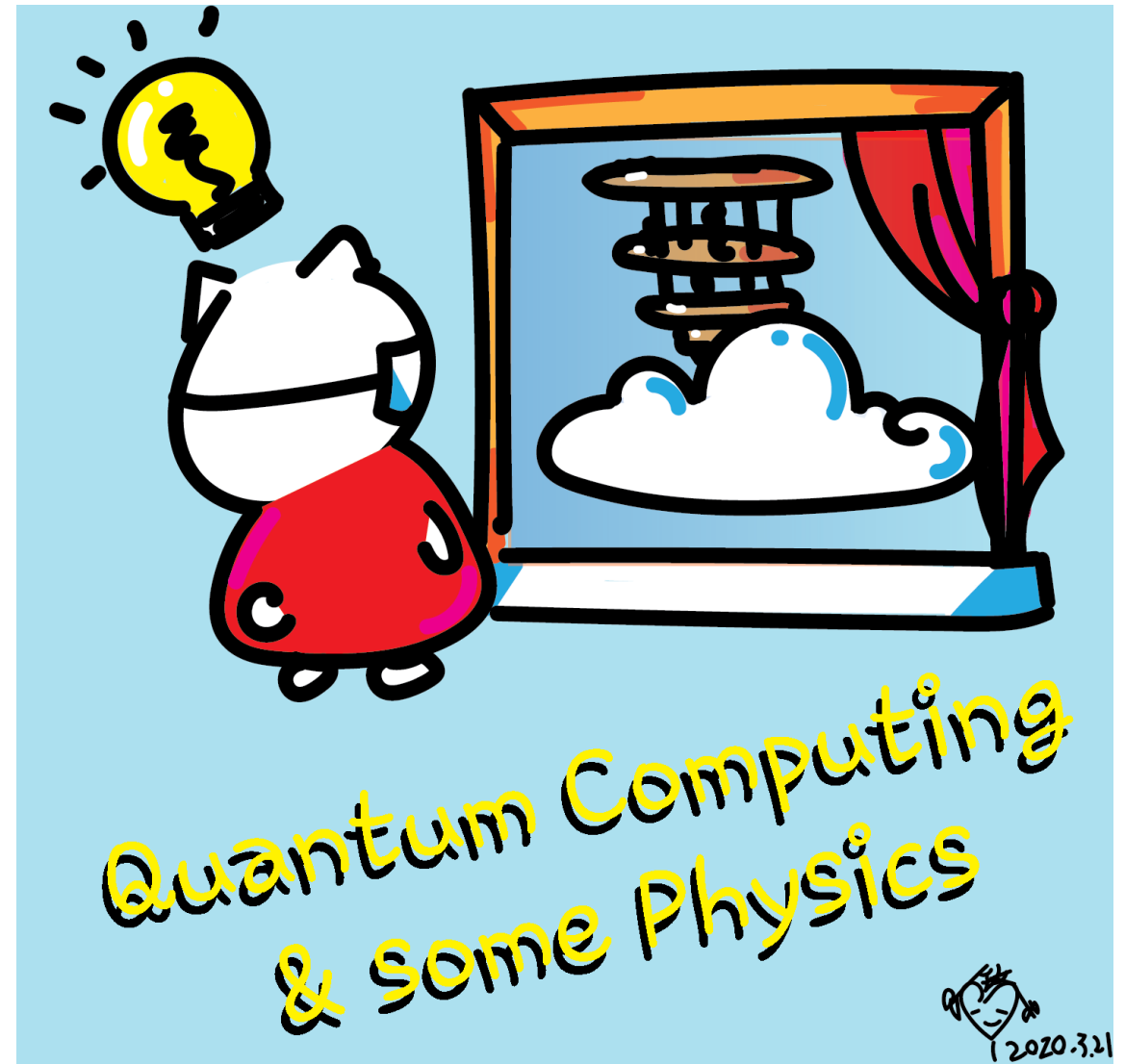
August 16, 2020

Hackaday, session 18

Other communities, session 10

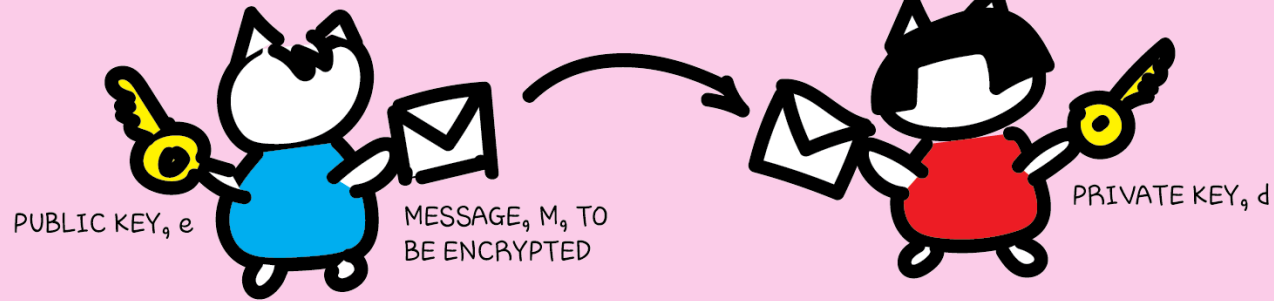
# Class structure

- [Comics on Hackaday – Quantum Computing through Comics](#) every Sun
- 30 mins – 1 hour every Sun, one concept (theory, hardware, programming), Q&A
- Contribute to Q# documentation  
<http://docs.microsoft.com/quantum>
- Coding through Quantum Katas  
<https://github.com/Microsoft/QuantumKatas/>
- Discuss in Hackaday project comments throughout the week
- Take notes



ENCRYPT MESSAGE USING  $e$ :  
 $C = M^e \text{ Mod } N$

DECRYPT CIPHER USING  $d$ :  
 $M = C^d \text{ Mod } N$



### THE RSA ENCRYPTION SCHEME

$$M^{ed} \text{ Mod } N = M$$

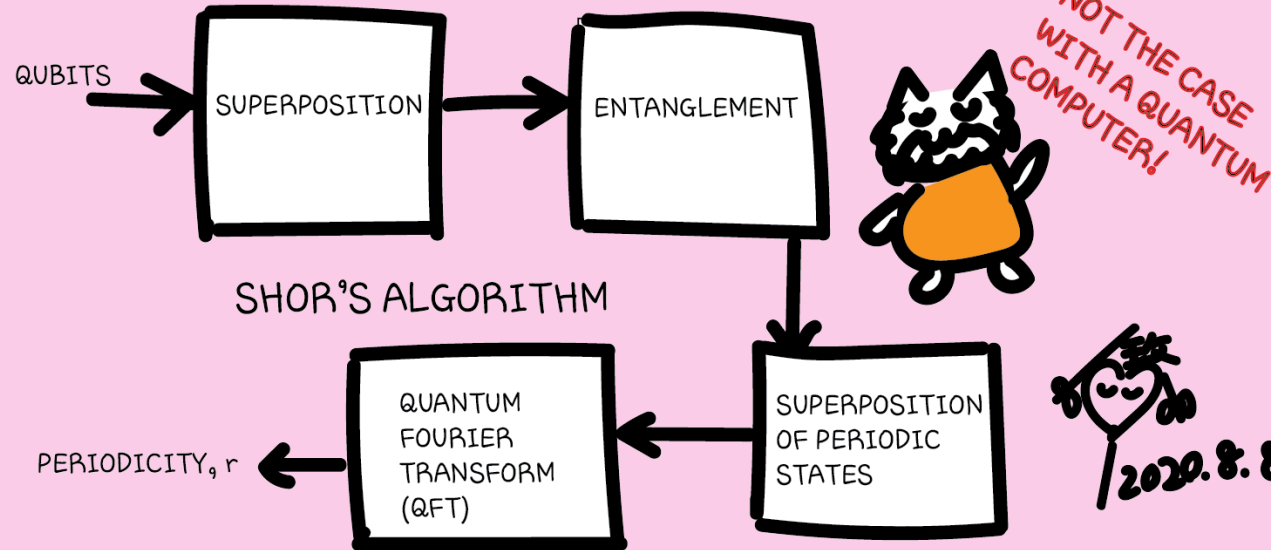
$$N = p * q$$

$$r = (p-1)(q-1)$$

$$e * d \text{ Mod } r = 1$$

PUBLIC:  $N, e$   
 PRIVATE:  $p, q, d, r$

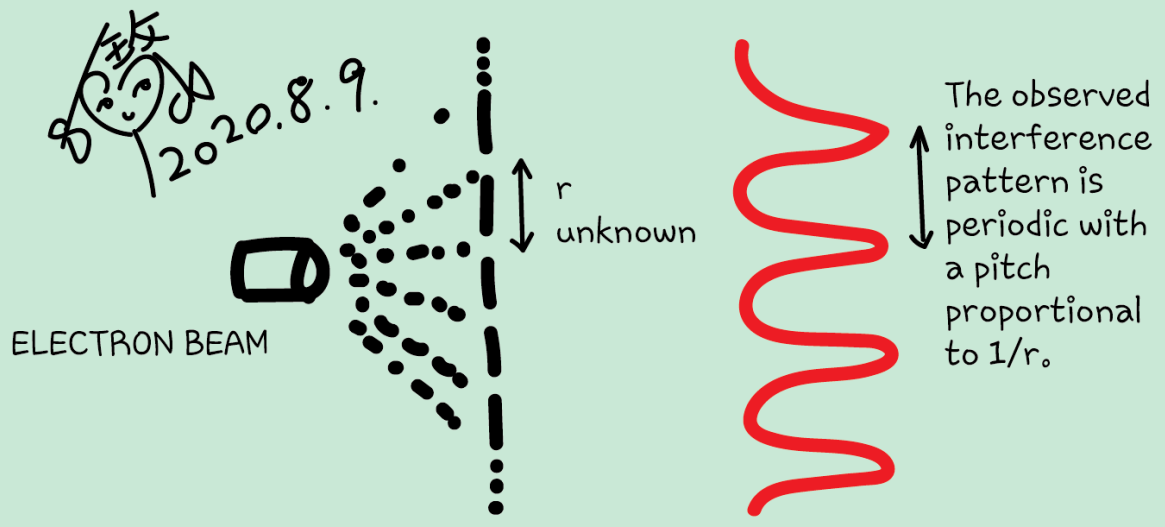
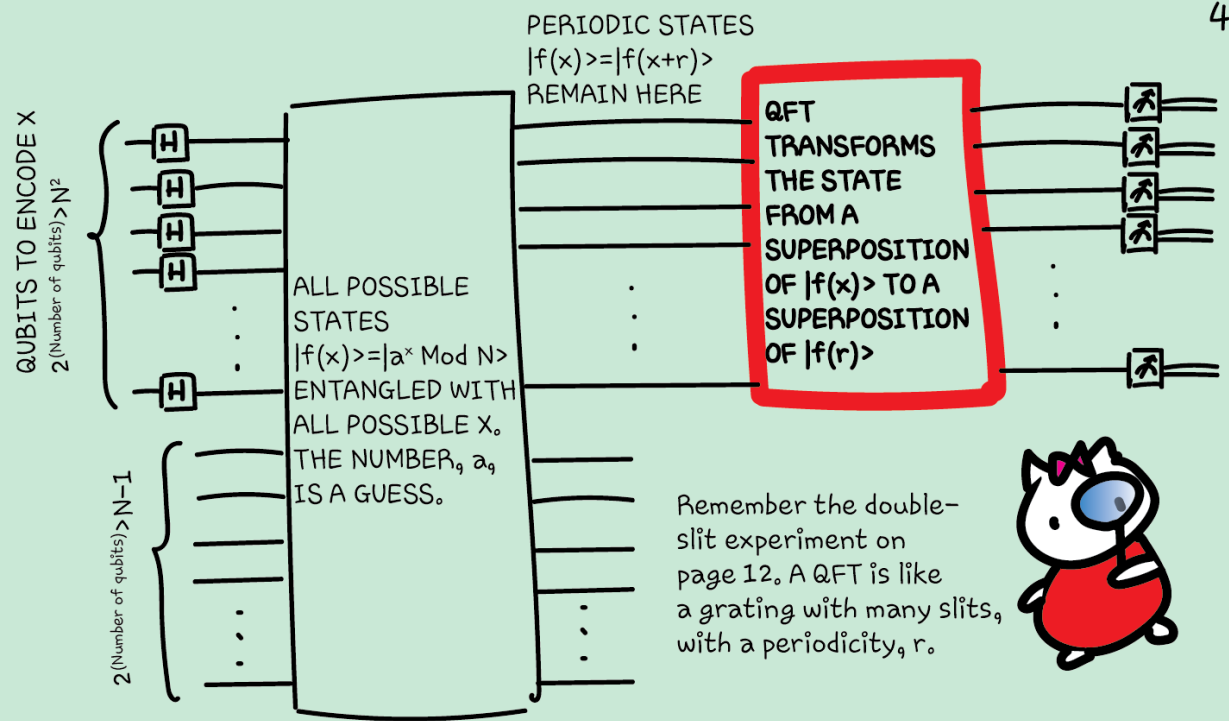
$N$  is really large - it is infeasible to factorize it classically to get  $p$  and  $q$ , thus,  $d$  and  $r$ .



### SHOR'S ALGORITHM

NOT THE CASE WITH A QUANTUM COMPUTER!

2020.8.8.



# RSA Numbers

- [https://en.wikipedia.org/wiki/RSA\\_numbers](https://en.wikipedia.org/wiki/RSA_numbers)
- RSA-100 has 100 decimal digits (330 bits). Its factorization was announced on April 1, 1991 by Arjen K. Lenstra. Reportedly, the factorization took a few days using the multiple-polynomial quadratic sieve algorithm on a MasPar parallel computer.
- RSA-100 = N =  
152260502792253336053561837813263742971806811496138068865790849458012296325895  
2897654000350692006139
- RSA-100 = p x q=  
37975227936943673922808872755445627854565536638199  
× 40094690950920881030683735292761468389214899724061
- Number of qubits needed  $\sim 659 + 329 = 988$   
(not considering error-correction qubits)

# Let's use a (much) smaller number

- $N = 35 = p * q$
- Number of qubits needed = 11, so that  $2^{11} = 2048 > N^2 = 35^2 = 1225$

$x$	$3^x \text{ Mod } 35$
0	1
1	3
2	9
3	27
4	11
5	33
6	29
7	17
8	16
9	13
10	4
11	12
12	1
13	3
14	9
15	27
16	11
17	33
18	29
19	17
20	16
21	13
22	4
23	12
24	1
25	3
26	9
.	.
.	.
.	.
2047	17

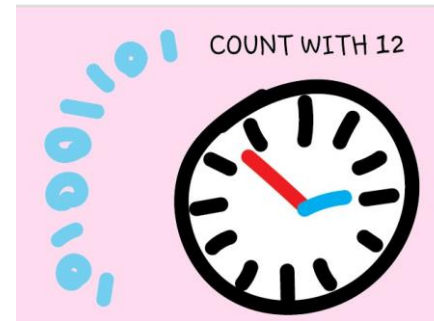
# Let's use a (much) smaller number

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Periodicity  $r = 12$

$$a^0 \text{ Mod } N = 3^0 \text{ Mod } 35 = a^r \text{ Mod } N = 3^{12} \text{ Mod } 35 = 1$$

$a^r - 1$  should be divisible by N



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7	17
8	16
9	13
10	4
11	12
12	1
13	3
14	9
15	27
16	11
17	33
18	29
19	17
20	16
21	13
22	4
23	12
24	1
25	3
26	9
.	.
.	.
.	.
2047	17

# Pure math

- Rearrange:  $(a^{r/2})^2 - 1$  should be divisible by  $N$
- $(a^{r/2} - 1)(a^{r/2} + 1)$  should be divisible by  $N$
- $r$  needs to be even,  $(a^{r/2} + 1)$  and  $(a^{r/2} - 1)$  are not individually divisible by  $N$
- $p$  divides  $(a^{r/2} - 1) = 728$ ,  $q$  divides  $(a^{r/2} + 1) = 730$
  
- $p = \text{GCD}(N, (a^{r/2} - 1)) = \text{GCD}(35, 728) = 7$
- $q = \text{GCD}(N, (a^{r/2} + 1)) = \text{GCD}(35, 730) = 5$

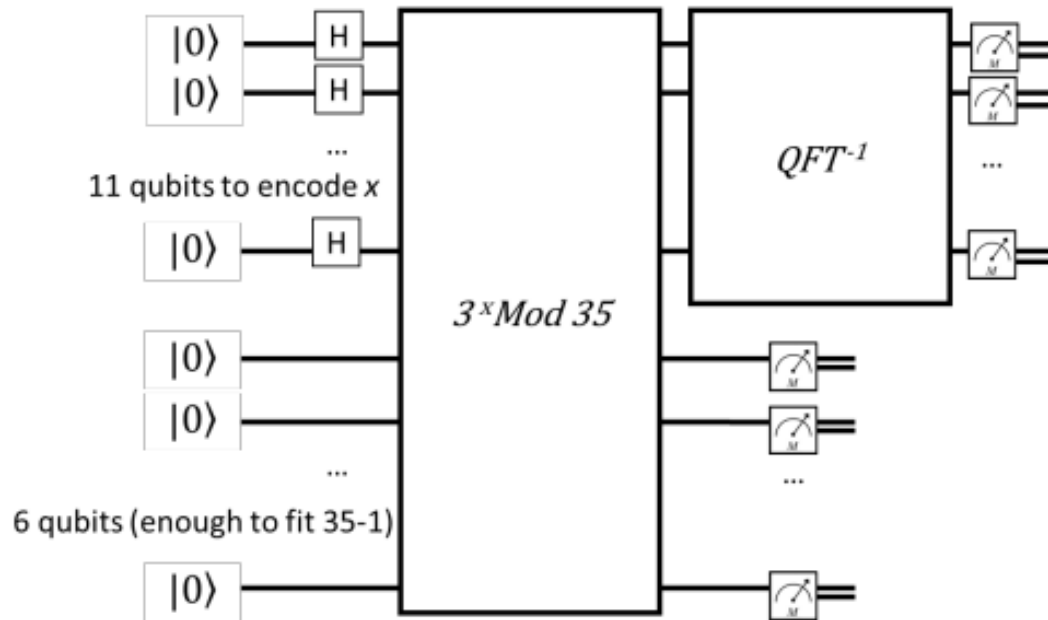


# Euclidean algorithm

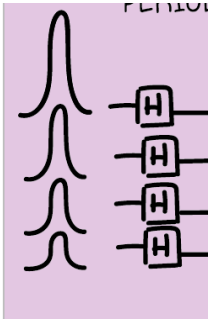
- $p = \text{GCD}(N, (a^{r/2} - 1)) = \text{GCD}(35, 728) = 7$
- $q = \text{GCD}(N, (a^{r/2} + 1)) = \text{GCD}(35, 730) = 5$
  
- $728 - 35 = 693$  divisible by 7
- $693 - 35 = 658$  divisible by 7
- ...
- $63 - 35 = 28$  divisible by 7

# Quantum part

- Hard step is finding  $r$  for large  $N$



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10	4
11	12
12	1
13	3
14	9
15	27
16	11
17	33
18	29
19	17
20	16
21	13
22	4
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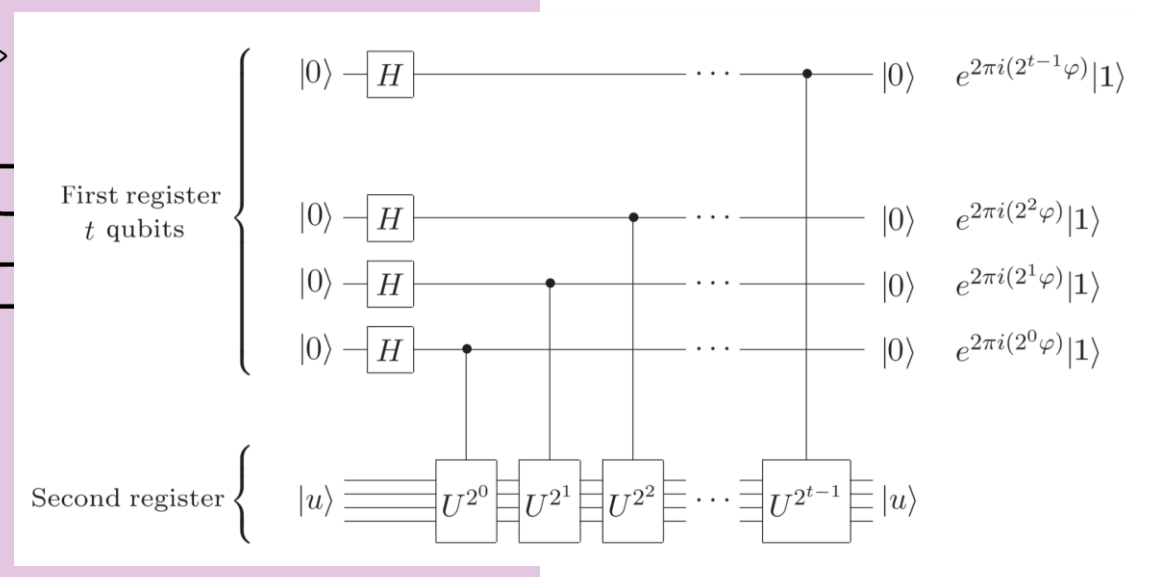
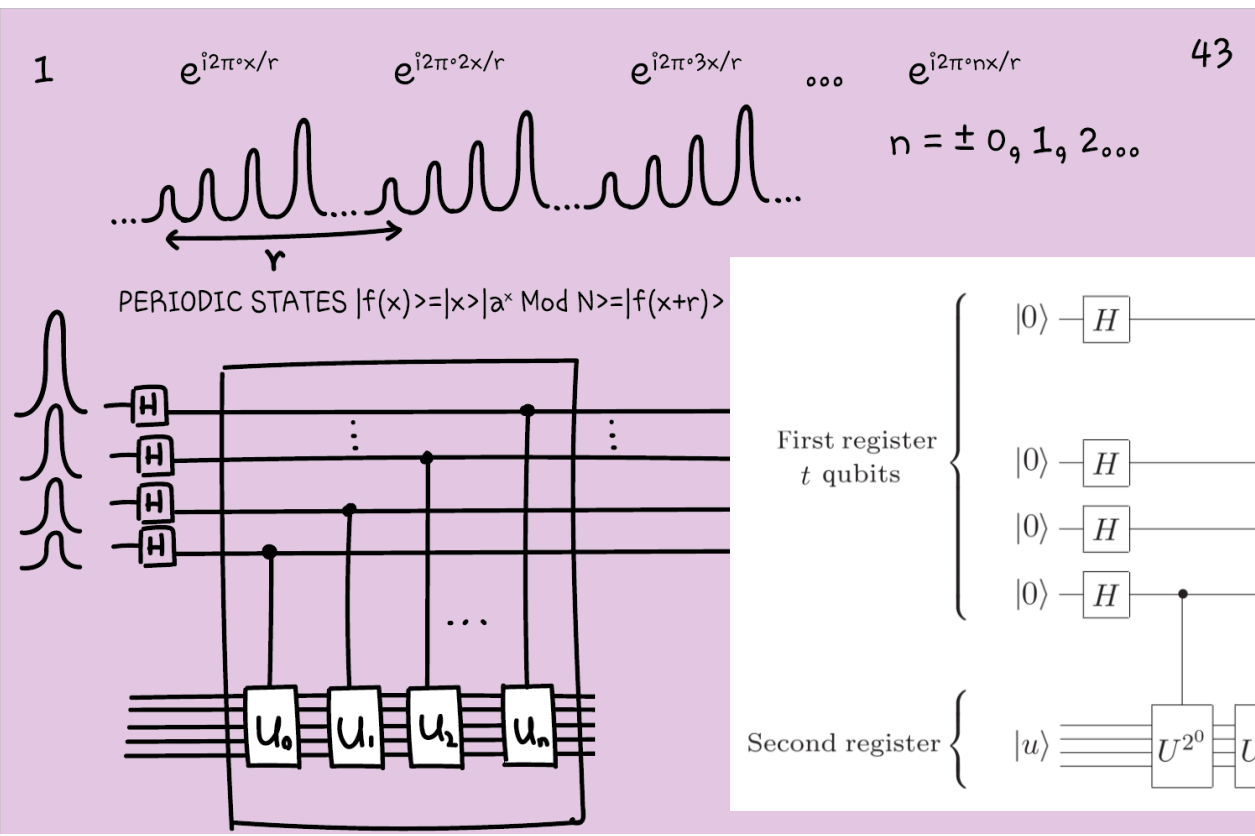
1  $e^{i2\pi x/r}$   $e^{i2\pi 2x/r}$   $e^{i2\pi 3x/r}$  ...  $e^{i2\pi nx/r}$  43

$n = \pm 0, 1, 2, \dots$

PERIODIC STATES  $|f(x)\rangle = |x\rangle |a^x \text{ Mod } N\rangle = |f(x+r)\rangle$

$$|0\rangle |a^0 \text{ Mod } N\rangle + |1\rangle |a^1 \text{ Mod } N\rangle + |2\rangle |a^2 \text{ Mod } N\rangle \dots |Q-1\rangle |a^{Q-1} \text{ Mod } N\rangle$$

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25	3
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$$|0\rangle|a^0 \text{ Mod } N\rangle + |1\rangle|a^1 \text{ Mod } N\rangle + |2\rangle|a^2 \text{ Mod } N\rangle \dots |Q-1\rangle|a^{Q-1} \text{ Mod } N\rangle$$

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19	17
20	16
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22	4
23	12
24	1
25	3
26	9
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2047	17

1. Quantum Computation and Quantum Information - 10th Anniversary Edition, Nielsen and Chuang  
 2. <https://github.com/Michaelvll/myQShor>  
 CS251 Quantum Information Science, 2018  
 @ ACM Honors Class, SJTU

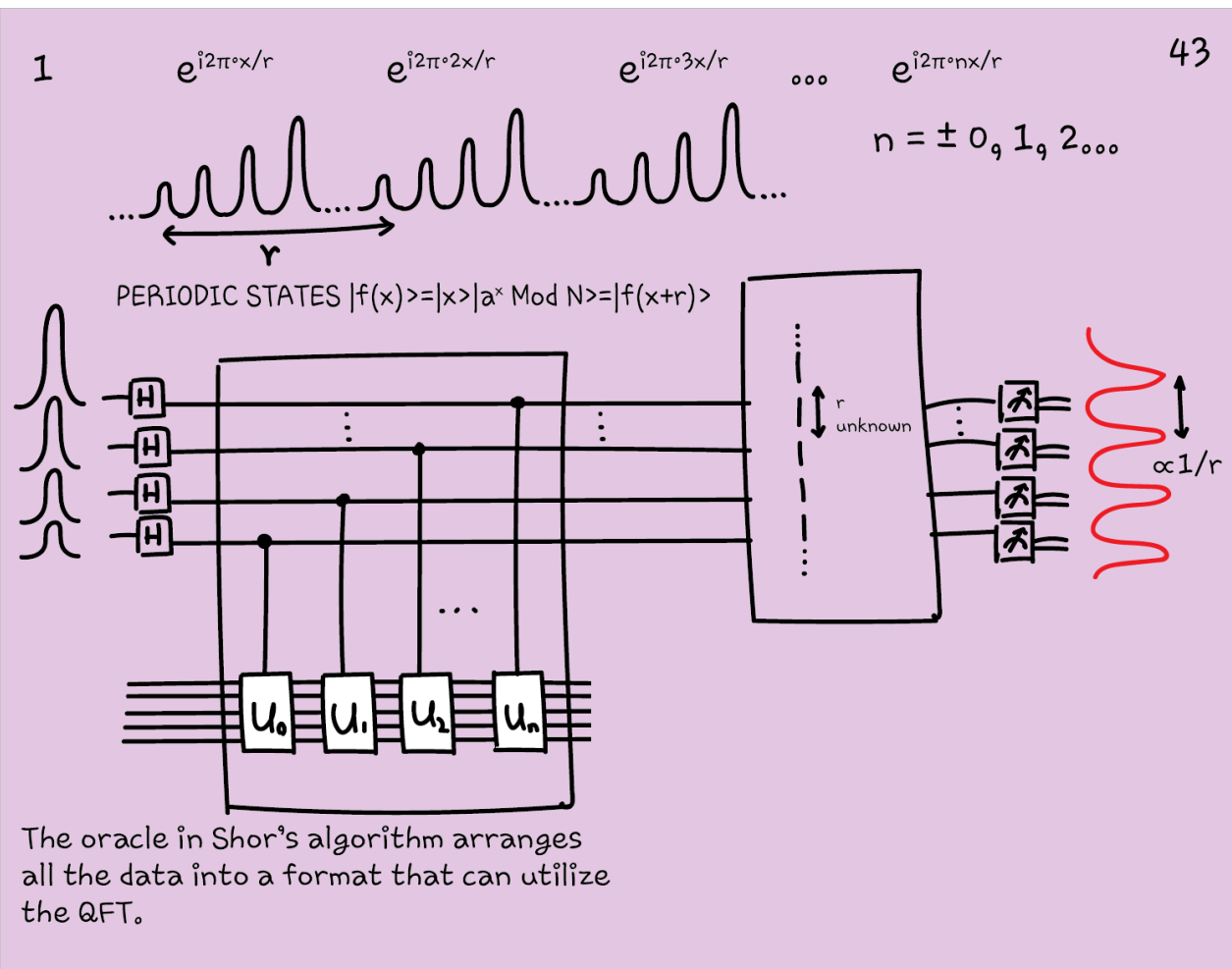
1  $e^{i2\pi x/r}$   $e^{i2\pi \cdot 2x/r}$   $e^{i2\pi \cdot 3x/r}$  ...  $e^{i2\pi \cdot nx/r}$  43

$n = \pm 0, 1, 2, \dots$

PERIODIC STATES  $|f(x)\rangle = |x\rangle |a^x \text{ Mod } N\rangle = |f(x+r)\rangle$

It is not required, but if one measured the output qubits here, one would obtain a periodic state with one of the amplitudes.

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1  $e^{i2\pi \cdot x/r}$   $e^{i2\pi \cdot 2x/r}$   $e^{i2\pi \cdot 3x/r}$  ...  $e^{i2\pi \cdot nx/r}$  43

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The oracle in Shor's algorithm arranges all the data into a format that can utilize the QFT.

### Box 5.1: Three qubit quantum Fourier transform

For concreteness it may help to look at the explicit circuit for the three qubit quantum Fourier transform:

Recall that  $S$  and  $T$  are the phase and  $\pi/8$  gates (see page xxiii). As a matrix the quantum Fourier transform in this instance may be written out explicitly, using  $\omega = e^{2\pi i/8} = \sqrt{i}$ , as

$$\frac{1}{\sqrt{8}} \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & \omega & \omega^2 & \omega^3 & \omega^4 & \omega^5 & \omega^6 & \omega^7 \\ 1 & \omega^2 & \omega^4 & \omega^6 & 1 & \omega^2 & \omega^4 & \omega^6 \\ 1 & \omega^3 & \omega^6 & \omega^1 & \omega^4 & \omega^7 & \omega^2 & \omega^5 \\ 1 & \omega^4 & 1 & \omega^4 & 1 & \omega^4 & 1 & \omega^4 \\ 1 & \omega^5 & \omega^2 & \omega^7 & \omega^4 & \omega^1 & \omega^6 & \omega^3 \\ 1 & \omega^6 & \omega^4 & \omega^2 & 1 & \omega^6 & \omega^4 & \omega^2 \\ 1 & \omega^7 & \omega^6 & \omega^5 & \omega^4 & \omega^3 & \omega^2 & \omega^1 \end{bmatrix} \quad (5.19)$$

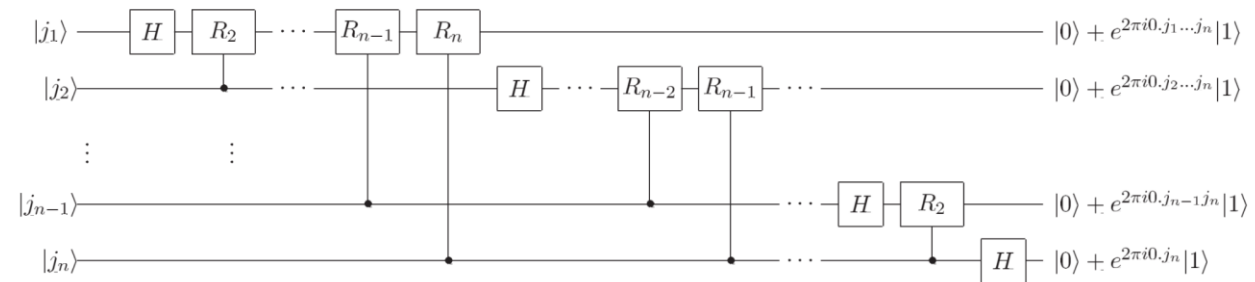


1  $e^{i2\pi \cdot x/r}$   $e^{i2\pi \cdot 2x/r}$   $e^{i2\pi \cdot 3x/r}$  ...  $e^{i2\pi \cdot nx/r}$  43

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The oracle in Shor's algorithm arranges all the data into a format that can utilize the QFT.



1.  $|0\rangle|0\rangle$  initial state
2.  $\rightarrow \frac{1}{\sqrt{2^t}} \sum_{x=0}^{2^t-1} |x\rangle|0\rangle$  create superposition
3.  $\rightarrow \frac{1}{\sqrt{2^t}} \sum_{x=0}^{2^t-1} |x\rangle|f(x)\rangle$  apply  $U$   
 $\approx \frac{1}{\sqrt{r2^t}} \sum_{\ell=0}^{r-1} \sum_{x=0}^{2^t-1} e^{2\pi i \ell x/r} |x\rangle |\hat{f}(\ell)\rangle$
4.  $\rightarrow \frac{1}{\sqrt{r}} \sum_{\ell=0}^{r-1} |\ell/r\rangle |\hat{f}(\ell)\rangle$  apply inverse Fourier transform to first register
5.  $\rightarrow \ell/r$  measure first register
6.  $\rightarrow r$  apply continued fractions algorithm

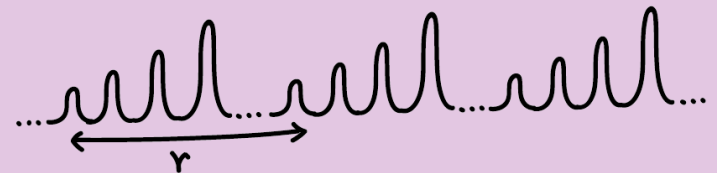
Quantum Computation and Quantum Information - 10th Anniversary Edition, Nielsen and Chuang

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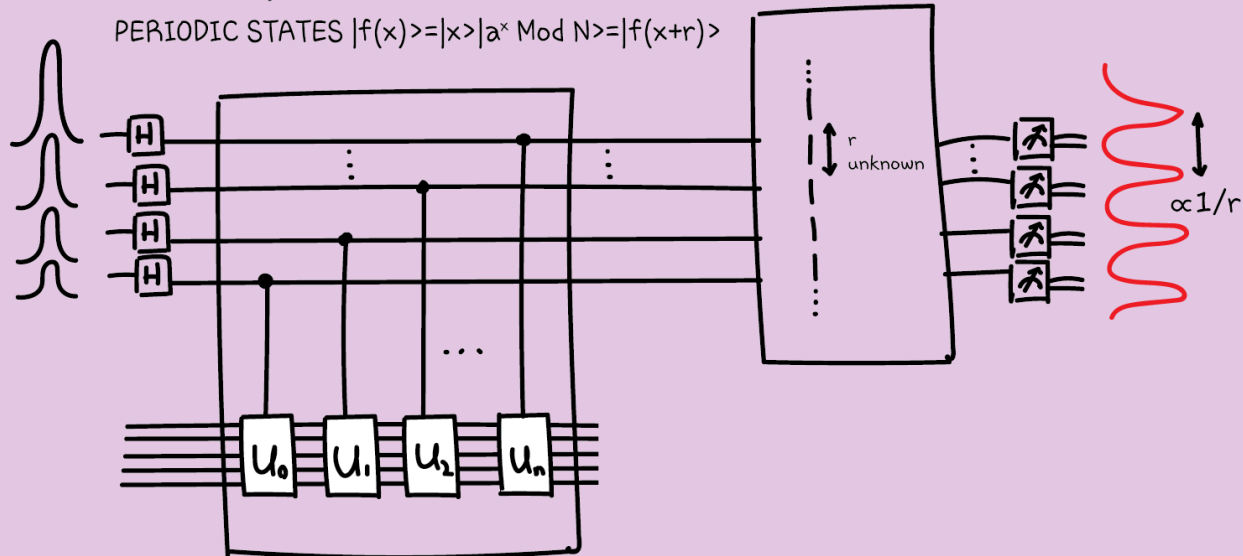
$$e^{i2\pi x/r} \quad e^{i2\pi \cdot 2x/r} \quad e^{i2\pi \cdot 3x/r} \quad \dots \quad e^{i2\pi \cdot nx/r}$$

43

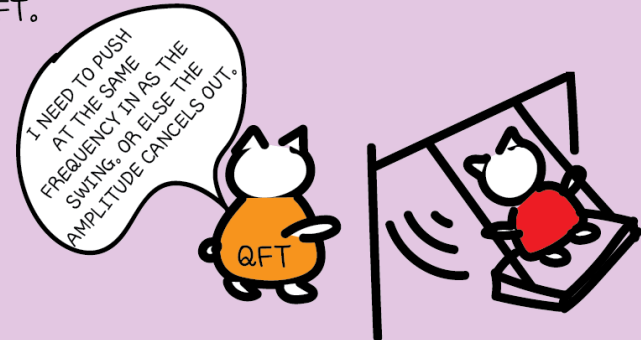
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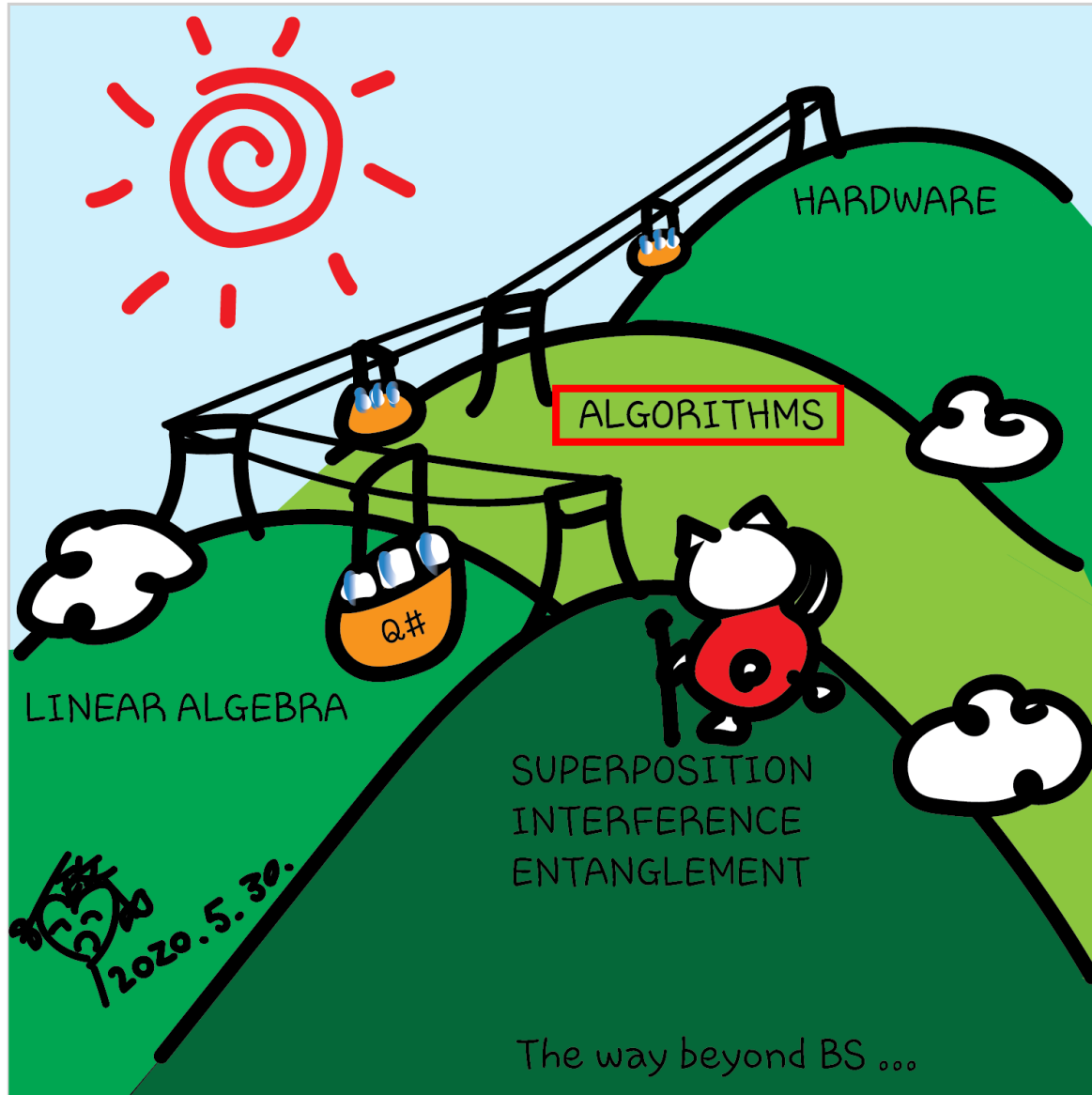
2020.8.16.

# Q# sample

- [microsoft/Quantum](#): Samples and tools to help get started with the Quantum Development Kit.
- Samples>algorithms>integer-factorization
- Numerics Library: Microsoft.Quantum.Arithmetic;
- QFT: Microsoft.Quantum.Canon

# Questions

- Post in chat or on Hackaday project  
<https://hackaday.io/project/168554-quantum-computing-through-comics>
- FAQ: Past Recordings on Hackaday project or my YouTube <https://www.youtube.com/c/DrKittyYeung>
- Next Sunday: quantum career Q&A session



The way beyond BS ...